A Lattice Compress and Forward Scheme

Natasha Devroye, *Assistant Professor, UIC* Yiwei Song, *Ph.D. candidate, UIC*



Gaussian networks



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Gaussian networks



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Structured codes for Gaussian networks



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Structured codes for Gaussian networks



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Structured codes for Gaussian networks





















• have: "decode the sum"









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• missing: "decode the sum"







• missing: "decode the sum"

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• missing: cooperation







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• AWGN relay channel





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``Cooperation"

Various links carry same message!



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- Decode and Forward (DF)
- Compress and Forward (CF)
- Amplify and Forward (AF)
- Quantize and Forward



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Lattice versions?



Central contribution

• AWGN relay channel





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• Decode and Forward (DF) → lattice codes achieve **DF** rate of random Gaussian codes [Song, Devroye ``List decoding for nested lattices

and applications to relay channels" Allerton 2010]



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- Decode and Forward (DF) → lattice codes achieve **DF** rate of random Gaussian codes [Song, Devroye ``List decoding for nested lattices and applications to relay channels" Allerton 2010]
 - **This paper** → lattice codes achieve **CF** rate of random Gaussian codes







• AWGN relay channel DF and CF schemes first considered in [Cover, El Gamal, 1979]



• DF extension to arbitrary # of relays and sources in [Xie, Kumar, 2004]





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 - Lattice DF single relay channel [Song, Devroye, Allerton 2010]

Outline - a lattice CF scheme

• Lattice notation

• Compress and forward review

• Lattice (X+Z1, X+Z2) Wyner-Ziv coding scheme

• Lattices achieve CF rate for AWGN relay

Lattice notation

- $\Lambda = \{\lambda = G \mathbf{i} : \mathbf{i} \in \mathbb{Z}^n\}, G$ the generator matrix
- lattice quantizer of Λ :

- $\mathbf{x} \mod \Lambda := \mathbf{x} Q(\mathbf{x})$
- fundamental region $\mathcal{V} := \{\mathbf{x} : Q(\mathbf{x}) = \mathbf{0}\}$ of volume V
- second moment per dimension of a uniform distribution over \mathcal{V} :

$$\sigma^2(\Lambda) := \frac{1}{V} \cdot \frac{1}{n} \int_{\mathcal{V}} ||\mathbf{x}||^2 d\mathbf{x}$$

 $Q(\mathbf{X}) = \arg\min_{\lambda \in \Lambda} ||\mathbf{X} - \lambda||$







Nested lattice codes

• Nested lattice pair : $\Lambda \subseteq \Lambda_c$ (Λ is Rogers-good and Poltyrev-good, Λ_c is Poltyrev-good)





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The (dithered) code book ★ C = {Λ_c ∩ V(Λ)} ★ is used to achieve the capacity of AWGN channel and used to achieve the quadratic Gaussian R(D) in [Erez, Litsyn, Zamir, Trans. IT, 2005] [Erez+Zamir, Trans. IT, 2004]





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• Coding rate:
$$R = \frac{1}{n} \log |\mathcal{C}| = \frac{1}{n} \log \frac{V(\Lambda)}{V(\Lambda_c)}$$
 arbitrary (# of \bigstar)




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• DF limited by need to decode at relay $R < I(X_1; Y_R | X_R)$



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$y_{R,1} \to \widehat{y_{R,1}}(i_1)$		



















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$$x_{R,1}(1)$$







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Compress + forward



ard		
$y_{R,1} \to \widehat{y_{R,1}}(i_1)$	$y_{R,2} \to \widehat{y_{R,2}}(i_2)$	
$x_{R,1}(1)$		



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R	$y_{R,1} \to \widehat{y_{R,1}}(i_1)$	$y_{R,2} \to \widehat{y_{R,2}}(i_2)$	
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Compress + forward $\begin{array}{c|c} & y_{R,1} \rightarrow \widehat{y_{R,1}}(i_1) & y_{R,2} \rightarrow \widehat{y_{R,2}}(i_2) \\ & x_{R,1}(1) & x_{R,2}(i_1) \end{array}$





Compress + forward $y_{R,1} \rightarrow \widehat{y_{R,1}(i_1)} \quad y_{R,2} \rightarrow \widehat{y_{R,2}(i_2)} \quad x_{R,1}(i_1) \quad x_{R,2}(i_1)$





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Combine

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Combine





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Key issue in CF: how to compress?

$$Y_R = X_1 + Z_R$$



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$$Y_R = X_1 + Z_R$$



direct link side-information

 $Y_2 = X_1 + Z_2$

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• demonstrate $(X + Z_1, X + Z_2)$ for completeness

$$\alpha_1 = \sqrt{1 - \frac{D}{N_1 + \frac{PN_2}{P + N_2}}}, \quad \alpha_2 = \frac{P}{P + N_2}$$

Theorem. The following rate-distortion function for the lossy compression of the source $X + Z_1$ subject to the reconstruction side-information $X + Z_2$ and squared error distortion metric may be achieved using lattice codes:

$$R(D) = \frac{1}{2} \log \left(\frac{\sigma_{X+Z_1|X+Z_2}^2}{D} \right), \qquad 0 \le D \le \sigma_{X+Z_1|X+Z_2}^2$$
$$= \frac{1}{2} \log \left(\frac{N_1 + \frac{PN_2}{P+N_2}}{D} \right), \qquad 0 \le D \le N_1 + \frac{PN_2}{P+N_2},$$

and 0 otherwise.

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Comparison with [Zamir, Erez, Shamai, 2002]



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Comment on α_1, α_2



Decoding (function $g(\cdot, \cdot)$)

 $\mathbf{X} + \mathbf{Z}_1 = \alpha_2(\mathbf{X} + \mathbf{Z}_2) + (1 - \alpha_2)\mathbf{X} + \mathbf{Z}_1 - \alpha_2\mathbf{Z}_2,$ \rightarrow choosing $\alpha_2 = \frac{P}{P+N_2}$, then $\mathbf{X} + \mathbf{Z_2} \perp (1 - \alpha_2)\mathbf{X} + \mathbf{Z_1} - \alpha_2\mathbf{Z_2}$

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 $\rightarrow \alpha_1 \equiv$ source coding MMSE coefficient

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 $\rightarrow \alpha_1 \equiv \text{source coding MMSE coefficient} \$ $\rightarrow \alpha_2 \equiv \text{channel coding MMSE coefficient} \$ Lose if do not pick these optimally

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R 1 2

A Lattice CF scheme

Theorem. For the three user Gaussian relay channel described by the input/output equations $Y_R = X_1 + N_R$ at the relay's receiver and $Y_2 = X_1 + X_R + N_2$ at the destination, with corresponding input and noise powers P_1, P_R, N_R, N_2 , the following rate may be achieved using lattice codes in a lattice Compress-and-Forward fashion:

$$R < \frac{1}{2} \log \left(1 + \frac{P_1}{N_2} + \frac{P_1 P_R}{P_1 N_R + P_1 N_2 + P_R N_R + N_R N_2} \right).$$

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same as that achieved by Gaussian codes in the CF scheme of [Cover, El Gamal, 1979]

Mimic all steps with lattice codes





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rate
$$R$$

 $v \leftrightarrow t_1, \Lambda_1 \subseteq \Lambda_{c1}, \sigma^2(\Lambda_1) = P_1$















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Encoding
Block b

$$x_{1,b}(w_b)$$

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$$y_{R,b} \rightarrow \widehat{y_{R,b}(i_b)}$$
send $x_{R,b}(i_{b-1})$

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$$x_{1} \in \hat{R} \quad \sigma^2(\Lambda) = N_R + \frac{P_1 N_2}{P_1 + N_2} + D$$

$$i \leftrightarrow t_R, \quad \Lambda_R \subseteq \Lambda_{cR}$$
rate $R' \quad \sigma^2(\Lambda_R) = P_R$

$$P_R$$

$$y'_{2,b-1} = X_{1,b-1} + Z_2$$

$$(-Y'_{2,b} = Y_{2,b} - X_{R,b}(i_{b-1}) = X_{1,b} + Z_2$$





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• decodes *i* from $Y_2 = X_1 + X_R + Z_2$ as long as $R' < \frac{1}{2} \log_2 \left(1 + \frac{P_R}{P_1 + N_2} \right)$





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• use $Y'_2 = Y_2 - X_R$ from previous block and X_R from current block to reconstruct \hat{Y}_R as in $(X + Z_R, X + Z_2)$ Wyner-Ziv



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• coherently combine Y'_2 and \hat{Y}_R to decode w, as long as $R < \frac{1}{2} \log \left(1 + \frac{P_1}{N_2} + \frac{P_1}{N_R + D} \right)$





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• note: pick $\alpha_1 = 1$ rather than source coding MMSE to render compression noise independent of all else





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 can this be combined with linearity of lattices to achieve higher rates in Gaussian networks?



• CF, DF lattices can achieve same rates as random Gaussian codebooks

 can this be combined with linearity of lattices to achieve higher rates in Gaussian networks?

• what is the capacity of relay channels, what are we stuck at?

Questions?

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